

DEFINITION 2021. $f/\sim = \uparrow \pi \circ f \circ \uparrow \pi^{-1}$ for every morphism f .

OBVIOUS 2022. $\text{Ob}(f/\sim) = (\text{Ob } f)/r$.

OBVIOUS 2023. $f/\sim = \langle \uparrow^{\text{FCD}} \pi \times^{(C)} \uparrow^{\text{FCD}} \pi \rangle f$ for every morphism f .

To define co-equalizers of morphisms f and g let \sim be is the smallest equivalence relation such that $fx = gx$.

LEMMA 2024. Let $f : X \rightarrow Y$ be a morphism of the category $\text{cont}(\mathcal{C})$ where \mathcal{C} is a concrete category (so $Wf = \uparrow \varphi$ for a **Rel**-morphism φ because f is principal) such that φ respects \sim . Factor it $\varphi = u \circ \pi$ where $u : \text{Ob}(X/\sim) \rightarrow \text{Ob } Y$ using properties of **Set**. Then u is a morphism of $\text{cont}(\mathcal{C})$ (that is a continuous function $X/\sim \rightarrow Y$).

PROOF. $f \circ X \circ f^{-1} \sqsubseteq Y$; $\uparrow u \circ \uparrow \pi \circ X \circ \uparrow \pi^{-1} \circ \uparrow u^{-1} \sqsubseteq Y$; $\uparrow u \in C(\uparrow \pi \circ X \circ \uparrow \pi^{-1}, Y) = C(X/\sim, Y)$. \square

THEOREM 2025. The following is a co-equalizer of parallel morphisms $f, g : A \rightarrow B$ of category $\text{cont}(\mathcal{C})$:

- the object $Y = f/\sim$;
- the morphism π considered as a morphism $B \rightarrow Y$.

PROOF. Let $z \circ f = z \circ g$ for some morphism z .

Let's prove $u \circ \pi = z$ for some $u : Y \rightarrow \text{Dst } z$. Really, as a morphism of **Set** it exists and is unique.

$\text{Src } z \in Y$. Thus $z = u \circ \pi$ for some u (by properties of **Set**). The function u is a $\text{cont}(\mathcal{C})$ -morphism because of the lemma above. It is unique by properties of **Set** (π obviously respects equivalence classes). \square

3. Rest

THEOREM 2026. The categories $\text{cont}(\mathcal{C})$ (for example in **Fcd** and **Rld**) are complete. **Fixme: Note that small complete category is a preorder!**

PROOF. They have products and equalizers. \square

THEOREM 2027. The categories $\text{cont}(\mathcal{C})$ (for example in **Fcd** and **Rld**) are co-complete.

PROOF. They have co-products and co-equalizers. \square

DEFINITION 2028. I call morphisms f and g of a category with embeddings *equivalent* ($f \sim g$) when there exist a morphism p such that $\text{Src } p \sqsubseteq \text{Src } f$, $\text{Src } p \sqsubseteq \text{Src } g$, $\text{Dst } p \sqsubseteq \text{Dst } f$, $\text{Dst } p \sqsubseteq \text{Dst } g$ and $\iota_{\text{Src } f, \text{Dst } f} p = f$ and $\iota_{\text{Src } g, \text{Dst } g} p = g$.

PROBLEM 2029. Find under which conditions:

- 1°. Equivalence of morphisms is an equivalence relation.
- 2°. Equivalence of morphisms is a congruence for our category.