

Chapter 11

Connectedness regarding functors and retracts

Definition 11.1. I will call *endoretracts* and *endofunctors* retracts and functors with the same source and destination. [TODO: Move above “continuity” chapter.]

11.1 Some lemmas

Lemma 11.2. If $\neg(A [f]^* B) \wedge A \cup B \in \text{dom } f \sqcup \text{im } f$ then f is closed on $\uparrow^U A$ for a functor $f \in \text{FCD}(U; U)$ for every sets U and $A, B \in \mathcal{P}U$.

Proof. Let $A \cup B \in \text{dom } f \sqcup \text{im } f$. $\neg(A [f]^* B) \Leftrightarrow \uparrow^U B \cap \langle f \rangle \uparrow^U A = 0^{\mathfrak{S}(U)} \Rightarrow (\text{dom } f \sqcup \text{im } f) \cap \uparrow^U B \cap \langle f \rangle \uparrow^U A = 0^{\mathfrak{S}(U)} \Rightarrow ((\text{dom } f \sqcup \text{im } f) \setminus \uparrow^U A) \cap \langle f \rangle \uparrow^U A = 0^{\mathfrak{S}(U)} \Leftrightarrow \langle f \rangle \uparrow^U A \subseteq \uparrow^U A$. \square

Corollary 11.3. If $\neg(A [f]^* B) \wedge A \cup B \in \text{dom } f \sqcup \text{im } f$ then f is closed on $\uparrow^U(A \setminus B)$ for a functor $f \in \text{FCD}(U; U)$ for every sets U and $A, B \in \mathcal{P}U$.

Proof. Let $\neg(A [f]^* B) \wedge A \cup B \in \text{dom } f \sqcup \text{im } f$. Then $\neg((A \setminus B) [f]^* B) \wedge (A \setminus B) \cup B \in \text{dom } f \sqcup \text{im } f$. \square

Lemma 11.4. If $\neg(A [f]^* B) \wedge A \cup B \in \text{dom } f \sqcup \text{im } f$ then $\neg(A [f^n]^* B)$ for every whole positive n .

Proof. Let $\neg(A [f]^* B) \wedge A \cup B \in \text{dom } f \sqcup \text{im } f$. From the above lemma $\langle f \rangle \uparrow^U A \subseteq \uparrow^U A$. $\uparrow^U B \cap \langle f \rangle \uparrow^U A = 0^{\mathfrak{S}(U)}$, consequently $\langle f \rangle \uparrow^U A \subseteq \uparrow^U(A \setminus B)$. Because (by the above corollary) f is closed on $\uparrow^U(A \setminus B)$, then $\langle f \rangle \langle f \rangle \uparrow^U A \subseteq \uparrow^U(A \setminus B)$, $\langle f \rangle \langle f \rangle \langle f \rangle \uparrow^U A \subseteq \uparrow^U(A \setminus B)$, etc. So $\langle f^n \rangle \uparrow^U A \subseteq \uparrow^U(A \setminus B)$, $\uparrow^U B \times \langle f^n \rangle \uparrow^U A = 0^{\mathfrak{S}(U)}$, $\neg(A [f^n]^* B)$. \square

11.2 Endomorphism series

Definition 11.5. $S_1(\mu) = \mu \sqcup \mu^2 \sqcup \mu^3 \sqcup \dots$ for an endomorphism μ of a precategory with countable join of morphisms (that is join defined for every countable set of morphisms).

Definition 11.6. $S(\mu) = \mu^0 \sqcup S_1(\mu) = \mu^0 \sqcup \mu \sqcup \mu^2 \sqcup \mu^3 \sqcup \dots$ where $\mu^0 = 1_{\text{Ob } \mu}$ (identity morphism for the object $\text{Ob } \mu$) where $\text{Ob } \mu$ is the object of endomorphism μ for an endomorphism μ of a precategory with countable join of morphisms.

I call S_1 and S *endomorphism series*.

We will consider the collection of all binary relations (on a set \mathcal{U}), as well as the collection of all functors and the collection of all retracts on a fixed set, as categories with single object \mathcal{U} and the identity morphisms $\text{id}_{\mathcal{U}}$, $\text{id}^{\text{FCD}(\mathcal{U})}$, $\text{id}^{\text{RLD}(\mathcal{U})}$.

Proposition 11.7. The relation $S(\mu)$ is transitive for the category of binary relations.