

(where $\langle p \rangle^*$ is defined in the introduction), recall that a funcoïd is uniquely determined by the values of its first component on sets. I will call such funcoïds *principal*. So funcoïds are a generalization of binary relations.

Composition of binary relations (i.e. of principal funcoïds) complies with the formulas:

$$\langle g \circ f \rangle^* = \langle g \rangle^* \circ \langle f \rangle^* \quad \text{and} \quad \langle (g \circ f)^{-1} \rangle^* = \langle f^{-1} \rangle^* \circ \langle g^{-1} \rangle^*.$$

By similar formulas we can define composition of every two funcoïds. Funcoïds with this composition form a category (*the category of funcoïds*).

Also funcoïds can be reversed (like reversal of X and Y in a binary relation) by the formula $(\alpha; \beta)^{-1} = (\beta; \alpha)$. In the particular case if μ is a proximity we have $\mu^{-1} = \mu$ because proximities are symmetric.

Funcoïds behave similarly to (multivalued) functions but acting on filters instead of acting on sets. Below these will be defined domain and image of a funcoïd (the domain and the image of a funcoïd are filters).

6.2. Basic definitions

DEFINITION 568. Let us call a *funcoïd* from a set A to a set B a quadruple $(A; B; \alpha; \beta)$ where $\alpha \in \mathfrak{F}(B)^{\mathfrak{F}(A)}$, $\beta \in \mathfrak{F}(A)^{\mathfrak{F}(B)}$ such that

$$\forall \mathcal{X} \in \mathfrak{F}(A), \mathcal{Y} \in \mathfrak{F}(B) : (\mathcal{Y} \neq \alpha \mathcal{X} \Leftrightarrow \mathcal{X} \neq \beta \mathcal{Y}).$$

Further we will assume that all funcoïds in consideration are small without mentioning it explicitly. **Fixme: It is superfluous.**

DEFINITION 569. *Source* and *destination* of every funcoïd $(A; B; \alpha; \beta)$ are defined as:

$$\text{Src}(A; B; \alpha; \beta) = A \quad \text{and} \quad \text{Dst}(A; B; \alpha; \beta) = B.$$

I will denote $\text{FCD}(A; B)$ the set of funcoïds from A to B .

I will denote FCD the set of all funcoïds (for small sets).

DEFINITION 570. $\langle (A; B; \alpha; \beta) \rangle \stackrel{\text{def}}{=} \alpha$ for a funcoïd $(A; B; \alpha; \beta)$.

DEFINITION 571. The *reverse* funcoïd $(A; B; \alpha; \beta)^{-1} = (B; A; \beta; \alpha)$ for a funcoïd $(A; B; \alpha; \beta)$.

NOTE 572. The reverse funcoïd is *not* an inverse in the sense of group theory or category theory.

PROPOSITION 573. If f is a funcoïd then f^{-1} is also a funcoïd.

PROOF. It follows from symmetry in the definition of funcoïd. \square

OBVIOUS 574. $(f^{-1})^{-1} = f$ for a funcoïd f .

DEFINITION 575. The relation $[f] \in \mathcal{P}(\mathfrak{F}(\text{Src } f) \times \mathfrak{F}(\text{Dst } f))$ is defined (for every funcoïd f and $\mathcal{X} \in \mathfrak{F}(\text{Src } f)$, $\mathcal{Y} \in \mathfrak{F}(\text{Dst } f)$) by the formula $\mathcal{X} [f] \mathcal{Y} \Leftrightarrow \mathcal{Y} \neq \langle f \rangle \mathcal{X}$.

OBVIOUS 576. $\mathcal{X} [f] \mathcal{Y} \Leftrightarrow \mathcal{Y} \neq \langle f \rangle \mathcal{X} \Leftrightarrow \mathcal{X} \neq \langle f^{-1} \rangle \mathcal{Y}$ for every funcoïd f and $\mathcal{X} \in \mathfrak{F}(\text{Src } f)$, $\mathcal{Y} \in \mathfrak{F}(\text{Dst } f)$.

OBVIOUS 577. $[f^{-1}] = [f]^{-1}$ for a funcoïd f .

THEOREM 578. Let A, B be sets.

- 1°. For given value of $\langle f \rangle \in \mathfrak{F}(B)^{\mathfrak{F}(A)}$ there exists no more than one funcoïd $f \in \text{FCD}(A; B)$.
- 2°. For given value of $[f] \in \mathcal{P}(\mathfrak{F}(A) \times \mathfrak{F}(B))$ there exists no more than one funcoïd $f \in \text{FCD}(A; B)$.