

that is

$$\begin{aligned} \bigsqcup \left\{ \frac{z \in \mathfrak{A}}{z \sqsubseteq a \wedge z \sqcap b = \perp} \right\} &= \\ \bigsqcup \left\{ \frac{z \in \text{atoms } a}{z \sqcap b = \perp} \right\} &= \\ \bigsqcup (\text{atoms } a \setminus \text{atoms } b). & \end{aligned}$$

□

3.5. Partially ordered categories

3.5.1. Definition.

DEFINITION 203. I will call a partially ordered (pre)category a (pre)category together with partial order \sqsubseteq on each of its Mor-sets with the additional requirement that

$$f_1 \sqsubseteq f_2 \wedge g_1 \sqsubseteq g_2 \Rightarrow g_1 \circ f_1 \sqsubseteq g_2 \circ f_2$$

for every morphisms f_1, g_1, f_2, g_2 such that $\text{Src } f_1 = \text{Src } f_2$ and $\text{Dst } f_1 = \text{Dst } f_2 = \text{Src } g_1 = \text{Src } g_2$ and $\text{Dst } g_1 = \text{Dst } g_2$.

3.5.2. Dagger categories.

DEFINITION 204. **FiXme:** In index I sometimes use *pre-category* and sometimes *precategory*. I will call a *dagger precategory* a precategory together with an involutive contravariant identity-on-objects prefunctor $x \mapsto x^\dagger$.

In other words, a dagger precategory is a precategory equipped with a function $x \mapsto x^\dagger$ on its set of morphisms which reverses the source and the destination and is subject to the following identities for every morphisms f and g :

- 1°. $f^{\dagger\dagger} = f$;
- 2°. $(g \circ f)^\dagger = f^\dagger \circ g^\dagger$.

DEFINITION 205. I will call a *dagger category* a category together with an involutive contravariant identity-on-objects functor $x \mapsto x^\dagger$.

In other words, a dagger category is a category equipped with a function $x \mapsto x^\dagger$ on its set of morphisms which reverses the source and the destination and is subject to the following identities for every morphisms f and g and object A :

- 1°. $f^{\dagger\dagger} = f$;
- 2°. $(g \circ f)^\dagger = f^\dagger \circ g^\dagger$;
- 3°. $(1_A)^\dagger = 1_A$.

THEOREM 206. If a category is a dagger precategory then it is a dagger category.

PROOF. We need to prove only that $(1_A)^\dagger = 1_A$. Really,

$$(1_A)^\dagger = (1_A)^\dagger \circ 1_A = (1_A)^\dagger \circ (1_A)^{\dagger\dagger} = ((1_A)^\dagger \circ 1_A)^\dagger = (1_A)^{\dagger\dagger} = 1_A.$$

□

For a partially ordered dagger (pre)category I will additionally require (for every morphisms f and g with the same source and destination)

$$f^\dagger \sqsubseteq g^\dagger \Leftrightarrow f \sqsubseteq g.$$

An example of dagger category is the category **Rel** whose objects are sets and whose morphisms are binary relations between these sets with usual composition of binary relations and with $f^\dagger = f^{-1}$.