

# Equalizers and co-Equalizers in Certain Categories

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## 1 Draft status

It is a rough draft. Errors are possible. Subscribe to my blog for further results.

<http://math.stackexchange.com/questions/540220/right-adjoint-of-forgetful-functor-from-top>

[TODO: Change notation  $\prod \rightarrow \prod^{(L)}$  .]

## 2 Categories with embeddings

**Note 1.** This section is not used below, it is just to feed your intuition.

The following generalizes the well known concept of embedding function  $A \hookrightarrow B$  for from a set  $A$  to a set  $B$  where  $A \subseteq B$ .

I will set that the unique morphism from an object  $A$  to an object  $B$  of a thin category is equal to the pair  $(A; B)$ .

**Definition 2.** A *category with embeddings of objects* is a dagger category with a preorder of the set of objects together with a functor  $\hookrightarrow$  (we will denote applying this functor to the object  $(A; B)$  as  $A \hookrightarrow B$ .) such that:

- $\hookrightarrow$  is an identity on objects.
- Every  $A \hookrightarrow B$  is a monomorphism.
- $(A \hookrightarrow B)^\dagger \circ (A \hookrightarrow B) = 1_A$ .

**Obvious 3.**  $A \hookrightarrow B$  is defined when  $(A; B)$  is a morphism of the preorder that is when  $A \sqsubseteq B$ .

**Obvious 4.**  $A \hookrightarrow B: A \rightarrow B$  when  $A \sqsubseteq B$ .

**Proposition 5.**  $A \hookrightarrow A = 1_A$ .

**Proof.** Because  $(A; A)$  is an identity morphism and  $\hookrightarrow$  preserves identities. □

**Proposition 6.**  $(B \hookrightarrow C) \circ (A \hookrightarrow B) = A \hookrightarrow C$  whenever  $A \sqsubseteq B \sqsubseteq C$ .

**Proof.**  $(B \hookrightarrow C) \circ (A \hookrightarrow B) = \hookrightarrow(B; C) \circ \hookrightarrow(A; B) = \hookrightarrow((B; C) \circ (A; B)) = \hookrightarrow(A; C) = A \hookrightarrow C$ . □

## 3 Categories under Rel

**Definition 7.** The **Rel**-morphism  $A \rightrightarrows B$  (*restriction-embedding*) is defined by the formula:  $A \rightrightarrows B = (A; B; \text{id}_{A \cap B})$ .

**Obvious 8.** If  $A \subseteq B$  then  $A \rightrightarrows B$  is an embedding  $A \hookrightarrow B = (A; B; \text{id}_A)$ .

**Obvious 9.** If  $A \supseteq B$  then  $A \rightrightarrows B = (A; B; \text{id}_B)$ .