

Connected functors and retracts*

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Abstract

Defined connectedness and connectivity for functors and retracts. Researched relationships of different kinds of connectedness. It is a part of my Algebraic General Topology research.

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1 Draft status

This article is a draft.

2 Prerequisites

For understanding this article prior reading of [1] is necessary.

3 Some lemmas

Lemma 1. If $\neg(A[f]B) \wedge A \cup B \supseteq \text{dom } f \cup \text{im } f$ then f is closed on A for a functor f and sets A and B .

Proof. $\neg(A[f]B) \Leftrightarrow B \cap \langle f \rangle A = \emptyset \Leftrightarrow (\text{dom } f \cup \text{im } f) \cap B \cap \langle f \rangle A = \emptyset \Rightarrow (\text{dom } f \cup \text{im } f \setminus A) \cap \langle f \rangle A = \emptyset \Leftrightarrow \langle f \rangle A \subseteq A$. \square

Corollary 2. If $\neg(A[f]B) \wedge A \cup B \supseteq \text{dom } f \cup \text{im } f$ then f is closed on $A \setminus B$ for a functor f and sets A and B .

Proof. Let $\neg(A[f]B) \wedge A \cup B \supseteq \text{dom } f \cup \text{im } f$. Then $\neg((A \setminus B)[f]B) \wedge (A \setminus B) \cup B \supseteq \text{dom } f \cup \text{im } f$. \square

Lemma 3. If $\neg(A[f]B) \wedge A \cup B \supseteq \text{dom } f \cup \text{im } f$ then $\neg(A[f^n]B)$ for any whole positive n .

Proof. Let $\neg(A[f]B) \wedge A \cup B \supseteq \text{dom } f \cup \text{im } f$. From the above proposition $\langle f \rangle A \subseteq A$. $B \cap \langle f \rangle A = \emptyset$, consequently $\langle f \rangle A \subseteq A \setminus B$. Because (by the above corollary) f is closed on $A \setminus B$, then $\langle f \rangle \langle f \rangle A \subseteq A \setminus B$, $\langle f \rangle \langle f \rangle \langle f \rangle A \subseteq A \setminus B$, etc. So $\langle f^n \rangle A \subseteq A \setminus B$, $B \cap \langle f^n \rangle A = \emptyset$, $\neg(A[f^n]B)$. \square

4 Endomorphism series

Definition 4. $S_1(\mu) \stackrel{\text{def}}{=} \mu \cup \mu^2 \cup \mu^3 \cup \dots$ for an endomorphism μ of a precategory with countable union of morphisms.

*. This document has been written using the GNU $\text{T}_{\text{E}}\text{X}_{\text{M}}\text{A}^{\text{C}}\text{S}$ text editor (see www.texmacs.org).

Definition 5. $S(\mu) \stackrel{\text{def}}{=} \mu^0 \cup S_1(\mu)$ where $\mu^0 \stackrel{\text{def}}{=} I_{\text{Ob } \mu}$ (identity morphism for the object $\text{Ob } \mu$) where $\text{Ob } \mu$ is the object of endomorphism μ for an endomorphism μ of a category with countable union of morphisms.

I call S_1 and S *endomorphism series*.

We will consider the collection of all binary relations (on a set \mathcal{U}), as well as the collection of all funcoids and the collection of all reloids, as categories with single object \mathcal{U} and the identity morphism $(=)$ or $(=)|_{\mathcal{U}}$.

So if μ is a binary relation or a funcoid or a reloid we have

$$S_1(\mu) = \mu \cup \mu^2 \cup \mu^3 \cup \dots \text{ and } S(\mu) = (=) \cup \mu \cup \mu^2 \cup \mu^3 \cup \dots$$

Proposition 6. $S(\mu)$ is transitive for the category of binary relations.

Proof.

$$\begin{aligned} S(\mu) \circ S(\mu) &= \mu^0 \circ S(\mu) \cup \mu \circ S(\mu) \cup \mu^2 \circ S(\mu) \cup \dots \\ &= (\mu^0 \cup \mu^1 \cup \mu^2 \cup \dots) \cup (\mu^1 \cup \mu^2 \cup \mu^3 \cup \dots) \cup (\mu^2 \cup \mu^3 \cup \mu^4 \cup \dots) \\ &= \mu^0 \cup \mu^1 \cup \mu^2 \cup \dots \\ &= S(\mu). \end{aligned}$$

□

5 Connectedness regarding binary relations

Before going to research connectedness for funcoids and reloids we will excuse into the basic special case of connectedness regarding binary relations.

Definition 7. A set A is called (*strongly*) *connected* regarding a binary relation μ when

$$\forall X, Y \in \mathcal{P}\mathcal{U} \setminus \{\emptyset\}: (X \cup Y \supseteq A \Rightarrow X[\mu]Y).$$

Definition 8. *Path* between two elements $a, b \in \mathcal{U}$ in a set A through binary relation μ is the finite sequence $x_0 \dots x_n$ where $x_0 = a$, $x_n = b$ for $n \in \mathbb{N}$ and $x_i(\mu \cap A \times A)x_{i+1}$ for any $i = 0, \dots, n - 1$. n is called *path length*.

Proposition 9. There exists path between any element $a \in \mathcal{U}$ and that element itself.

Proof. It is the path consisting of one vertex (of length 0). □

Proposition 10. There is a path from element a to element b in a set A through a binary relation μ iff $a(S(\mu \cap A \times A))b$ (that is $(a, b) \in S(\mu \cap A \times A)$).

Proof.

Direct implication. If exists a path from a to b , then $\{b\} \subseteq \langle (\mu \cap A \times A)^n \rangle \{a\}$ where n is the path length. Consequently $\{b\} \subseteq \langle S(\mu \cap A \times A) \rangle \{a\}$; $a(S(\mu \cap A \times A))b$.

Reverse implication. If $a(S(\mu \cap A \times A))b$ then exists $n \in \mathbb{N}$ such that $a(\mu \cap A \times A)^n b$. By definition of composition of binary relations this means that there exist finite sequence $x_0 \dots x_n$ where $x_0 = a$, $x_n = b$ for $n \in \mathbb{N}$ and $x_i(\mu \cap A \times A)x_{i+1}$ for any $i = 0, \dots, n - 1$. That is there is path from a to b . □

Theorem 11. The following statements are equivalent for a relation μ and a set A :

1. For any $a, b \in A$ there is a path between a and b in A through μ .
2. $S(\mu \cap A \times A) \supseteq A \times A$. [TODO: \supseteq can be replaced with $=$.]
3. A is connected regarding μ .

Proof.

- (1) \Rightarrow (2). Let for any $a, b \in A$ there is a path between a and b in A through μ . Then $a(S(\mu \cap A \times A))b$ for any $a, b \in A$. It is possible only when $S(\mu \cap A \times A) \supseteq A \times A$.
- (2) \Rightarrow (1). For any two vertices a and b we have $a(S(\mu \cap A \times A))b$. So (by the previous theorem) for any two vertices a and b exist path from a to b .
- (2) \Rightarrow (3). Suppose that $\neg(X[\mu \cap A \times A]Y)$ for some $X, Y \in \mathcal{P}U \setminus \{\emptyset\}$ such that $X \cup Y \supseteq A$. Then by a lemma $\neg(X[(\mu \cap A \times A)^n]Y)$ for any $n \in \mathbb{N}$. Consequently $\neg(X[S(\mu \cap A \times A)]Y)$. So $S(\mu \cap A \times A) \neq A \times A$.
- (3) \Rightarrow (2). If $\langle S(\mu \cap A \times A) \rangle \{v\} = A$ for every vertex v then $S(\mu \cap A \times A) = A \times A$. Consider the remaining case when $V \stackrel{\text{def}}{=} \langle S(\mu \cap A \times A) \rangle \{v\} \subset A$ for some vertex v . Let $W = A \setminus V$. Then $V \cup W \supseteq A$ and consequently $V \cap A \cup W \cap A = (V \cup W) \cap A = A$ consequently $V \cap A[\mu]W \cap A$ that is $V[\mu \cap A \times A]W$ that is $\langle \mu \cap A \times A \rangle V \cap W \neq \emptyset$. This is impossible because $\langle \mu \cap A \times A \rangle V = \langle \mu \cap A \times A \rangle \langle S(\mu \cap A \times A) \rangle V = \langle S_1(\mu \cap A \times A) \rangle V \subseteq \langle S(\mu \cap A \times A) \rangle V = V$. \square

Corollary 12. A set A is connected regarding a relation μ iff it is connected regarding $\mu \cap A \times A$. [TODO: Generalize this for funcoids and reloids.]

Definition 13. A *connected component* of a set A regarding a binary relation F is a maximal connected subset of A .

Theorem 14. The set A is partitioned into connected components (regarding any binary relation F).

Proof. Consider the binary relation $a \sim b \Leftrightarrow a(S(F))b \wedge b(S(F))a$. \sim is a symmetric, reflexive, and transitive relation. So all points of A are partitioned into a collection of sets Q . Obviously each component is (strongly) connected. If a set $R \subseteq A$ is greater than one of that connected components A then it contains a point $b \in B$ where B is some other connected component. Consequently R is disconnected. \square

Proposition 15. A set is connected (regarding a binary relation) iff it has one connected component.

Proof. Direct implication is obvious. Reverse is proved by contradiction. \square

6 Connectedness regarding funcoids and reloids

Definition 16. $S_1^*(\mu) = \bigcap^{\mathcal{F}} \{S_1(M) \mid M \in \text{up } \mu\}$ for a reloid μ .

Definition 17. *Connectivity reloid* for a reloid μ is defined as follows:

$$S^*(\mu) = \bigcap^{\mathcal{F}} \{S(M) \mid M \in \text{up } \mu\}.$$

Remark 18. Do not mess the word *connectivity* with the word *connectedness* which means being connected.¹

Proposition 19. $S^*(\mu) = (=) \cup S_1^*(\mu)$ for any reloid μ .

Proof. Follows from a theorem about distributivity of \cup regarding $\bigcap^{\mathcal{F}}$. \square

Proposition 20. $S^*(\mu) = S(\mu)$ if μ is a binary relation.

Proof. $S^*(\mu) = \bigcap^{\mathcal{F}} \{S(\mu)\} = S(\mu)$. \square

1. In some math literature these two words are used interchangeably.

Definition 21. A filter \mathcal{A} is called *connected* regarding a reloid μ when $S^*(\mu) \supseteq \mathcal{A} \times \mathcal{A}$.

Definition 22. A filter \mathcal{A} is called *connected* regarding a funcoid μ when

$$\forall X, Y \in \mathcal{P}\mathcal{U} \setminus \{\emptyset\}: (X \cup Y \supseteq \mathcal{A} \Rightarrow X[\mu]Y).$$

Theorem 23. A filter \mathcal{A} is connected regarding a funcoid μ iff

$$\forall \mathcal{X}, \mathcal{Y} \in \mathcal{F} \setminus \{\emptyset\}: (\mathcal{X} \cup \mathcal{Y} \supseteq \mathcal{A} \Rightarrow \mathcal{X}[\mu]\mathcal{Y}).$$

Proof.

Reverse implication. Obvious.

Direct implication. Let $\mathcal{X}, \mathcal{Y} \in \mathcal{F}$, $\mathcal{X} \cup \mathcal{Y} \supseteq \mathcal{A}$. Then $X \cup Y \supseteq \mathcal{A}$ for any $X \in \text{up}\mathcal{X}$, $Y \in \text{up}\mathcal{Y}$. Consequently $\forall X \in \text{up}\mathcal{X}, Y \in \text{up}\mathcal{Y}: X[\mu]Y$; $\mathcal{X}[\mu]\mathcal{Y}$. \square

Theorem 24. A filter \mathcal{A} is connected regarding a funcoid μ iff \mathcal{A} is connected for every binary relation $F \supseteq \mu$.

Proof.

Direct implication. Obvious.

Reverse implication. Let \mathcal{A} is not connected regarding μ , that is $\neg(X[\mu]Y)$ for some $X, Y \in \mathcal{P}\mathcal{U} \setminus \{\emptyset\}$ such that $X \cup Y \supseteq \mathcal{A}$. We can assume that $X \cap Y = \emptyset$. (Otherwise replace Y with $Y \setminus X$ or X with $X \setminus Y$; if both $Y \setminus X$ and $X \setminus Y$ are empty then $X = Y = \mathcal{A}$ what is impossible.) Also we can assume $X \cup Y = \mathcal{A}$. Consider the relation $F = (X \times X) \cup (Y \times \mathcal{A})$. For any atomic filter a

$$\langle F \rangle a = \left(\begin{cases} X & \text{if } a \subseteq X \\ \mathcal{A} & \text{if } a \not\subseteq X \end{cases} \right) \supseteq \langle \mu \cap \mathcal{A} \times \mathcal{A} \rangle a$$

So $\mu \cap \mathcal{A} \times \mathcal{A} \subseteq F$. Consequently $\mu \cap \mathcal{A} \times \mathcal{A} \subseteq F$ for some $A \in \text{up}\mathcal{A}$. So

$$F_2 = F \cup (\mathcal{U} \times \mathcal{U} \setminus A \times A) \supseteq \mu$$

Obviously $\neg(X[F]Y)$ and consequently $\neg(X[F_2]Y)$. So \mathcal{A} is not connected regarding F_2 . \square

Theorem 25. The following statements are equivalent for a binary relation F and a filter \mathcal{A} :

1. \mathcal{A} is connected regarding F as a funcoid.
2. \mathcal{A} is connected regarding F as a reloid.

Proof. From the previous section. \square

Theorem 26. A filter \mathcal{A} is connected regarding a reloid f iff it is connected regarding every $F \in \text{up}f$.

Proof.

Direct implication. Obvious.

Reverse implication. F is connected iff $S(F) = F^0 \cup F^1 \cup F^2 \cup \dots \supseteq \mathcal{A} \times \mathcal{A}$.

$$S^*(f) = \bigcap^{\mathcal{F}} \{S(F) \mid F \in \text{up}f\} \supseteq \bigcap^{\mathcal{F}} \{\mathcal{A} \times \mathcal{A} \mid F \in \text{up}f\} = \mathcal{A} \times \mathcal{A}. \quad \square$$

Theorem 27. A filter \mathcal{A} is connected regarding a reloid f iff it is connected regarding the funcoid (FCD) f .

Proof. f is connected iff every element of $\text{up}f$ is connected. But

$$\text{up}f = \{F \mid F \in \mathcal{P}\mathcal{U} \wedge F \supseteq (\text{FCD})f\}.$$

From the theorem 24 follows that these are equivalent. \square

7 Algebraic properties of S and S^*

Theorem 28. $S^*(S^*(f)) = S^*(f)$ for any reloid f .

Proof. $S^*(S^*(f)) = \bigcap^{\mathcal{F}} \{S(R) \mid R \in \text{up } S^*(f)\} \subseteq \bigcap^{\mathcal{F}} \{S(R) \mid R \in \{S(F) \mid F \in \text{up } f\}\} = \bigcap^{\mathcal{F}} \{S(S(F)) \mid F \in \text{up } f\} = \bigcap^{\mathcal{F}} \{S(F) \mid F \in \text{up } f\} = S^*(f)$.

So $S^*(S^*(f)) \subseteq S^*(f)$. $S^*(S^*(f)) \supseteq S^*(f)$ is obvious. \square

Corollary 29. $S^*(S(f)) = S(S^*(f)) = S^*(f)$ for any reloid f .

Proof. Obviously $S^*(S(f)) \supseteq S^*(f)$ and $S(S^*(f)) \supseteq S^*(f)$.

But $S^*(S(f)) \subseteq S^*(S^*(f)) = S^*(f)$ and $S(S^*(f)) \subseteq S^*(S^*(f)) = S^*(f)$. \square

Conjecture 30. $S(S(f)) = S(f)$ for

1. any reloid f ;
2. any funcoid f .

Conjecture 31. For any reloid f

1. $S(f) \circ S(f) = S(f)$;
2. $S^*(f) \circ S^*(f) = S^*(f)$;
3. $S(f) \circ S^*(f) = S^*(f) \circ S(f) = S^*(f)$.

Conjecture 32. $S(f) \circ S(f) = S(f)$ for any funcoid f .

Bibliography

- [1] Victor Porton. Funcoids and reloids. At <http://www.mathematics21.org/binaries/funcoids-reloids.pdf>.